

# Strategies for Parallelizing the Solution of Rational Matrix Equations

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# Main goal

Compare different strategies for parallelizing the solution of numerical problems

- Levels of granularity: Application, linear algebra operations, numerical kernel.
- Programming models: message-passing, shared memory.
- Difficulty of programming. Transparency to the programmer.
- Sequential and parallel libraries and tools: ScaLAPACK, LAPACK, BLAS, MPI.

# Outline

## 1 Rational Matrix Equation

## 2 Sequential algorithm

## 3 Parallel algorithms

- Multithreading algorithm (MT)
- Message-passing multiprocess algorithm (MP)
- Message-passing two processes algorithm (2P)
- Two-processes hybrid algorithm (HB2P)
- Multiprocess hybrid algorithm (HBMP)

## 4 Conclusions



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# Rational Matrix Equation (RME)

$$X = Q + LX^{-1}L^T$$

$Q \in \mathbb{R}^{n \times n}$  is symmetric positive definite,  $L \in \mathbb{R}^{n \times n}$  is nonsingular

$X_+ \in \mathbb{R}^{n \times n}$ : positive definite symmetric solution.

- **Applications:** Analysis of stationary Gaussian reciprocal processes.
- **Examples:**
  - Steady-state distribution of temperature along a heated beam subjected to random loads.
  - Ship surveillance problem. Prediction of its trajectory.



# Structure-preserving Doubling Algorithm (SDA)

- Based on the solution of a Discrete-time Algebraic Riccati equation.
- Iterative method.
  - Quadratic convergence rate.
  - Fair numerical stability.
  - Cubic computational cost on  $n$ .
- It can be efficiently parallelized.

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## Sequential algorithm

$$L_0 = \widehat{L}, \quad Q_0 = \widehat{Q} + \widehat{P}, \quad P_0 = 0$$

**for**  $i = 0, 1, 2, \dots$

$$Q_i - P_i = C_i^T C_i \quad \text{DPOTRF } (n^3/3)$$

$$A_C = L_i^T C_i^{-1} \quad \text{DTRSM } (n^3)$$

$$C_A = C_i^{-T} L_i^T \quad \text{DTRSM } (n^3)$$

$$Q_{i+1} = Q_i - C_A^T C_A \quad \text{DGEMM } (n^3)$$

$$P_{i+1} = P_i + A_C A_C^T \quad \text{DGEMM } (n^3)$$

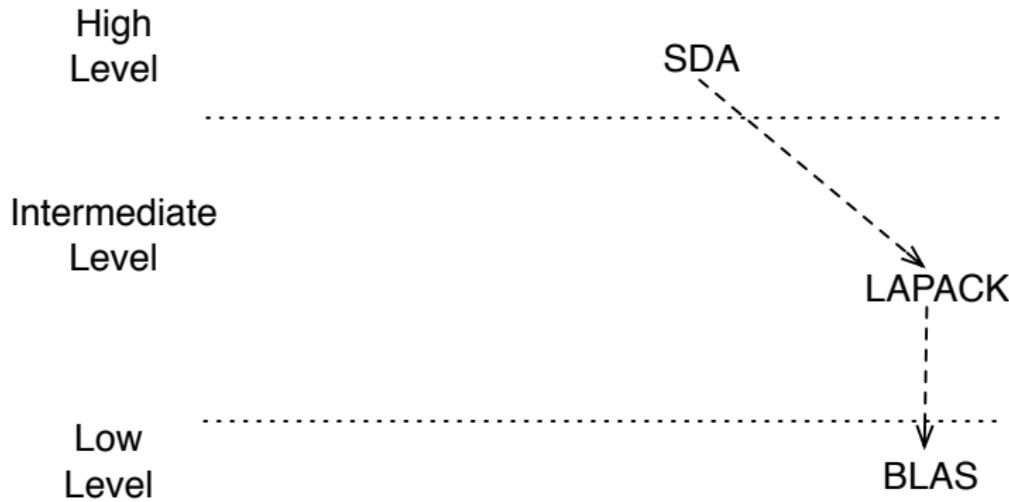
$$L_{i+1} = A_C C_A \quad \text{DGEMM } (2n^3)$$

**until** convergence

**Total cost**  $\approx 6.3n^3 flops$



# Levels of granularity



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# Parallelizing Strategies

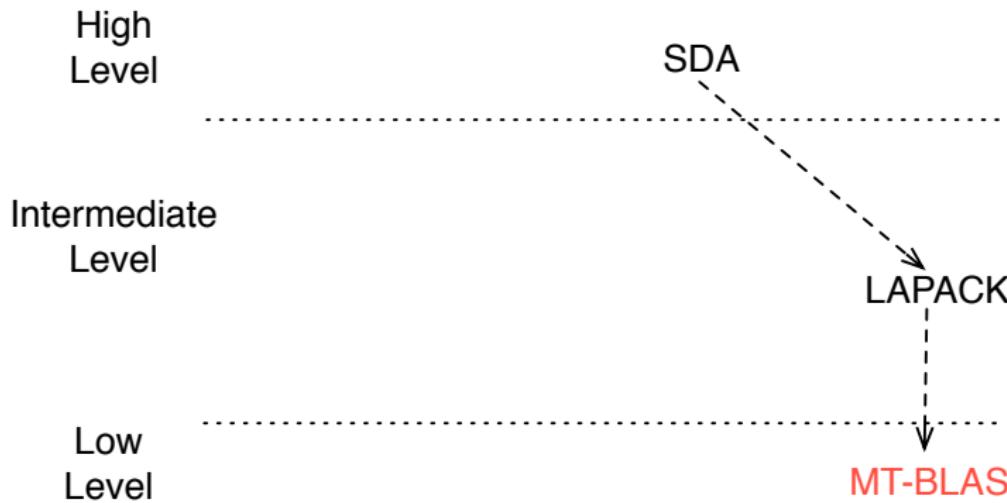
Exploit three levels of granularity:

- **High level:** Overlap steps of the SDA Algorithm.
- **Medium level:** Linear Algebra routines: Linear system solving, Cholesky factorization, etc.
- **Low level:** Numerical kernels: matrix products, matrix-vector products, etc.



# Multithreading algorithm (MT)

Lazy programmer



# Hardware environment: set

## SGI Altix 350

- DSM multiprocessor.
- 16 Itanium2 (1.5GHz.) processors (6 MB L3 cache).
- 32 GBytes of DSM.
- cc-NUMA architecture.
  - SGI NUMAlink interconnect.
  - Ring topology.
  - 8 dual-processor nodes.

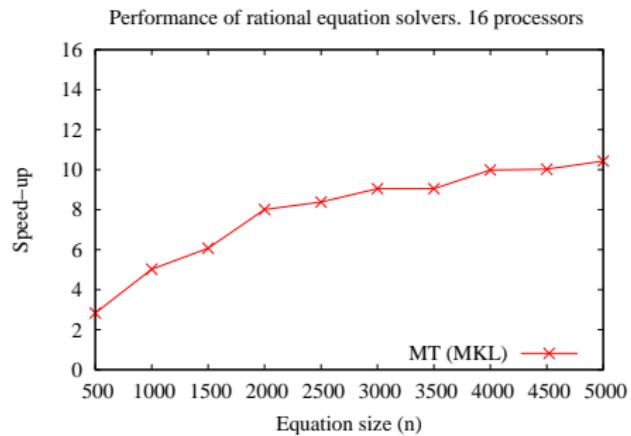
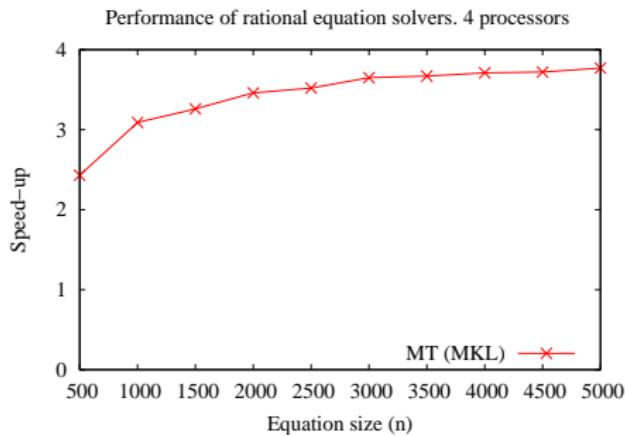


# Software environment

- Programming language: Fortran.
- Libraries and tools:
  - SGI ScaLAPACK for Distributed Shared Memory: SDSM.
  - SGI implementation of MPI.
  - Two multithreaded versions of LAPACK+BLAS:
    - SGI Scientific Computing Software Library: SCSL.
    - Intel Math Kernel Library: MKL.



# Experimental results



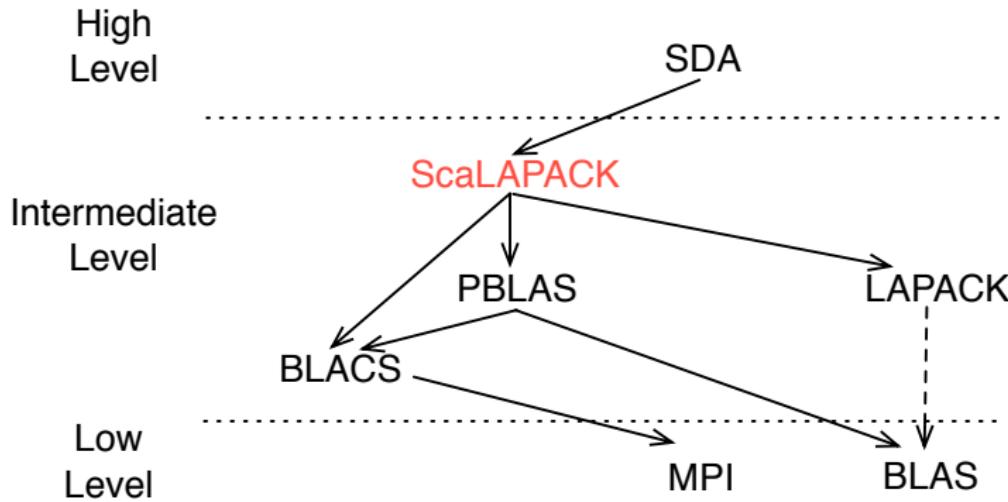
## Comments

- Can only be used in shared memory multiprocessors.
- Low performance when many processors are used. Possible causes:
  - Fine-grain parallelism.
  - Unique copy of the data in memory. Cost of the data transference to maintain the coherence of the data.

Could we improve the performance with a little more work?



# Message-passing multiprocess algorithm (MP)



# Multiprocess algorithm

## Data distribution

$$L_0 = \widehat{L}, \quad Q_0 = \widehat{Q} + \widehat{P}, \quad P_0 = 0$$

**for**  $i = 0, 1, 2, \dots$

$$Q_i - P_i = C_i^T C_i$$

PDPOTRF

$$A_C = L_i^T C_i^{-1}$$

PDTRSM

$$C_A = C_i^{-T} L_i^T$$

PDTRSM

$$Q_{i+1} = Q_i - C_A^T C_A$$

PDGEMM

$$P_{i+1} = P_i + A_C A_C^T$$

PDGEMM

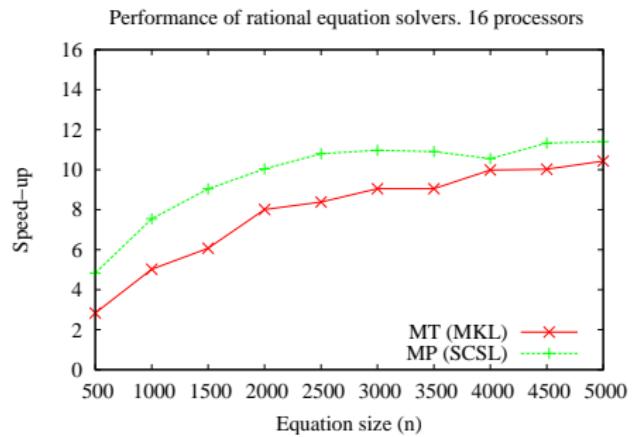
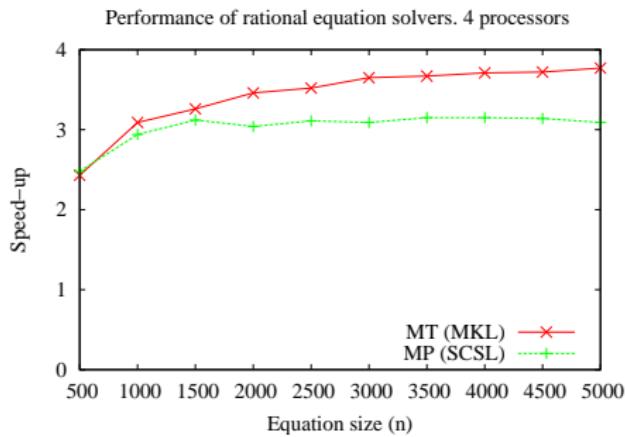
$$L_{i+1} = A_C C_A$$

PDGEMM

**until** convergence



# Experimental results



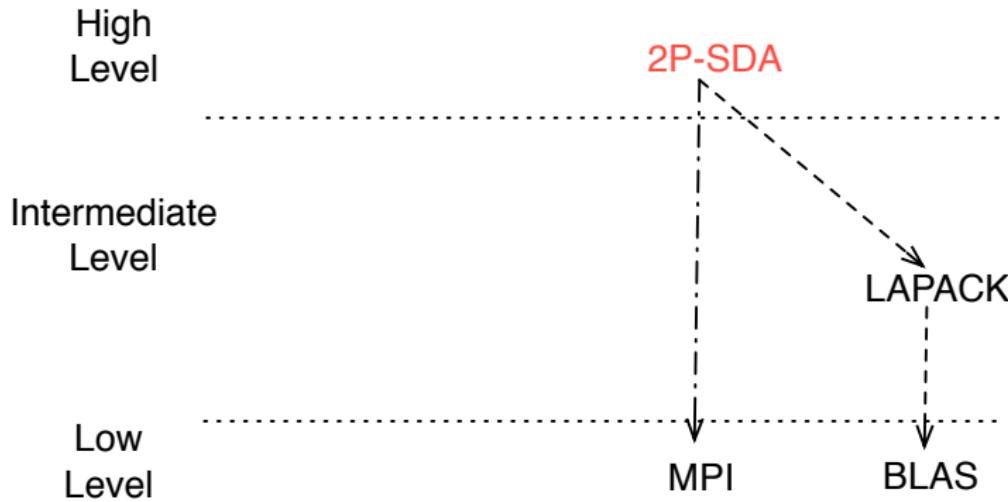
## Comments

- Can be used on shared and distributed memory multiprocessors.
- Worse performance than MT algorithm with few processors.
- Improves the performance of MT algorithm with more processors.
  - Coarser-grain parallelism than MT algorithm.
  - Distributed data among "local" memories. Larger locality to access data ⇒ Less data transference cost.

Could we improve the performance with a little more work? Again



# Message-passing two processes algorithm (2P)



# Two processes algorithm

$$L_0 = \widehat{L}, \quad Q_0 = \widehat{Q} + \widehat{P}, \quad P_0 = 0$$

**for**  $i = 0, 1, 2, \dots$

1.  $Q_i - P_i = C_i^T C_i$  DPOTRF ( $n^3/3$ )
2.  $A_C = L_i^T C_i^{-1}$  DTRSM ( $n^3$ )
3.  $C_A = C_i^{-T} L_i^T$  DTRSM ( $n^3$ )
4.  $Q_{i+1} = Q_i - C_A^T C_A$  DGEMM ( $n^3$ )
5.  $P_{i+1} = P_i + A_C A_C^T$  DGEMM ( $n^3$ )
- 6.1.  $L_{i+1} = A_C C_A$  DGEMM ( $n^3$ )
- 6.2.  $L_{i+1} = A_C C_A$  DGEMM ( $n^3$ )

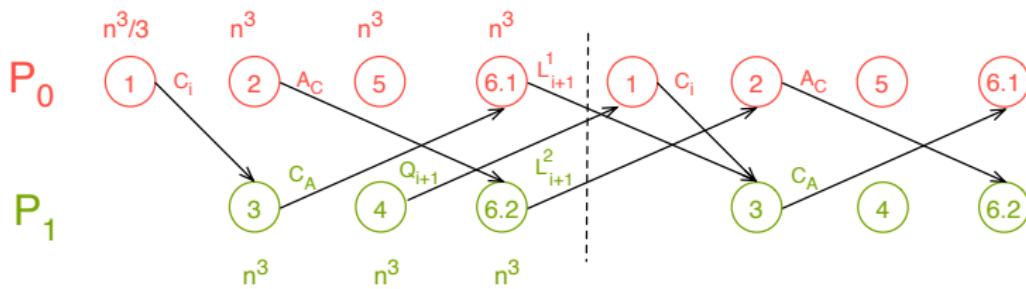
**until** convergence

$$P_0 \text{ cost} \approx 3.3n^3 \text{ flops}$$

$$P_1 \text{ cost} \approx 3n^3 \text{ flops}$$



# Dependencies and Communications



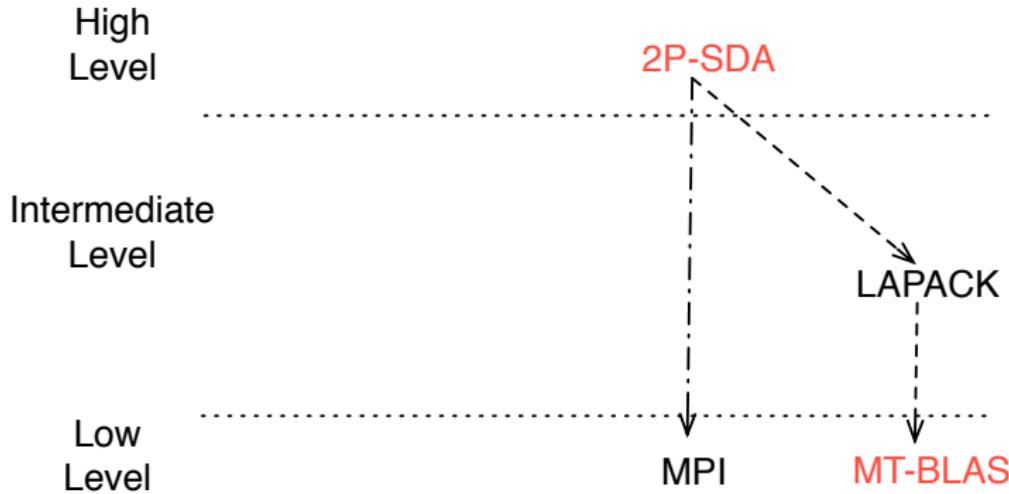
## Comments

- Can be used on shared and distributed memory multiprocessors.
- Exploits the highest level parallelism.
- Uses only two MPI processes.

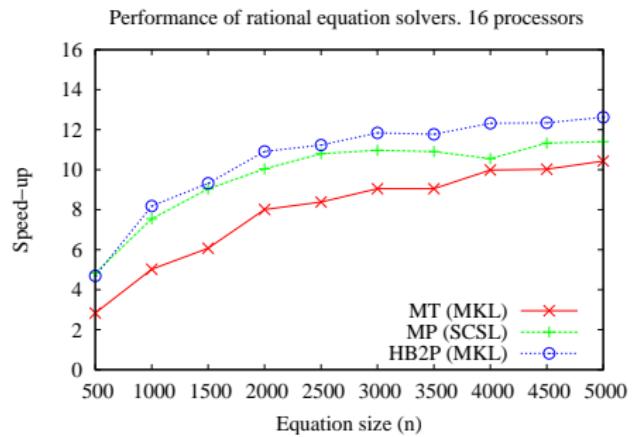
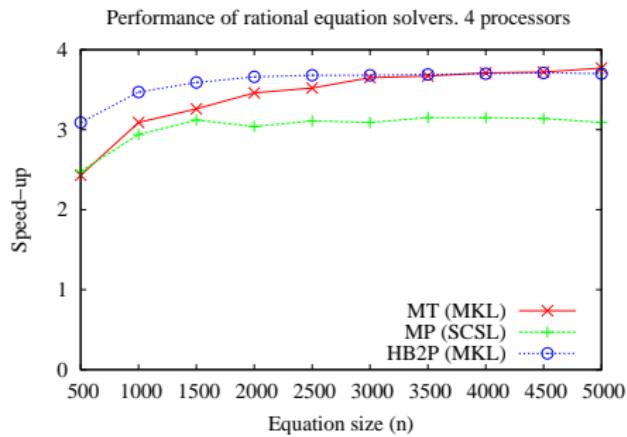
Could we take profit of more than two processors?



# Two-processes hybrid algorithm (HB2P)

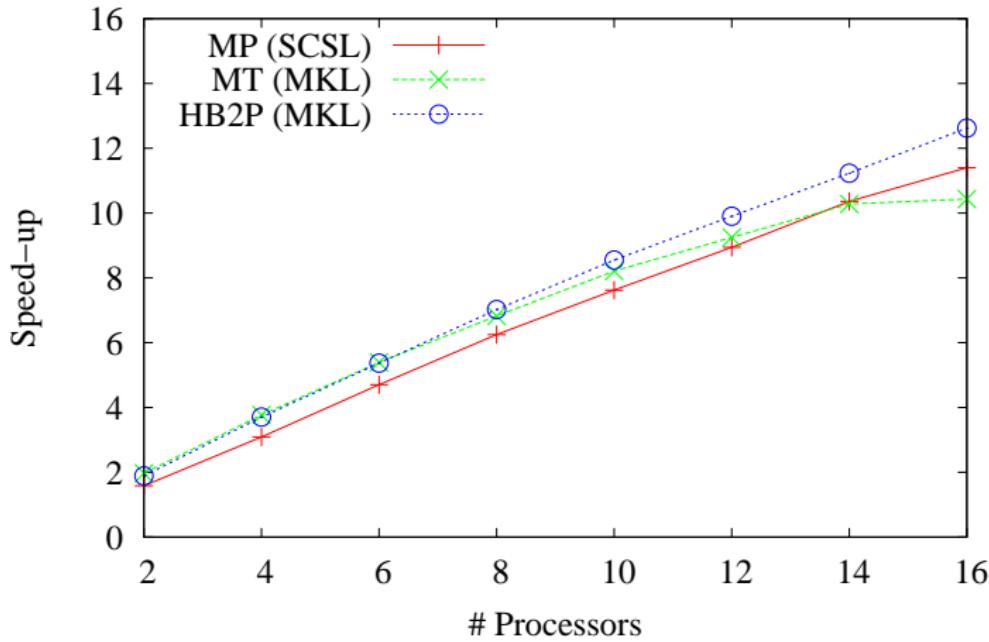


# Experimental results



# Experimental results

Performance of rational equation solvers.  $n=5000$



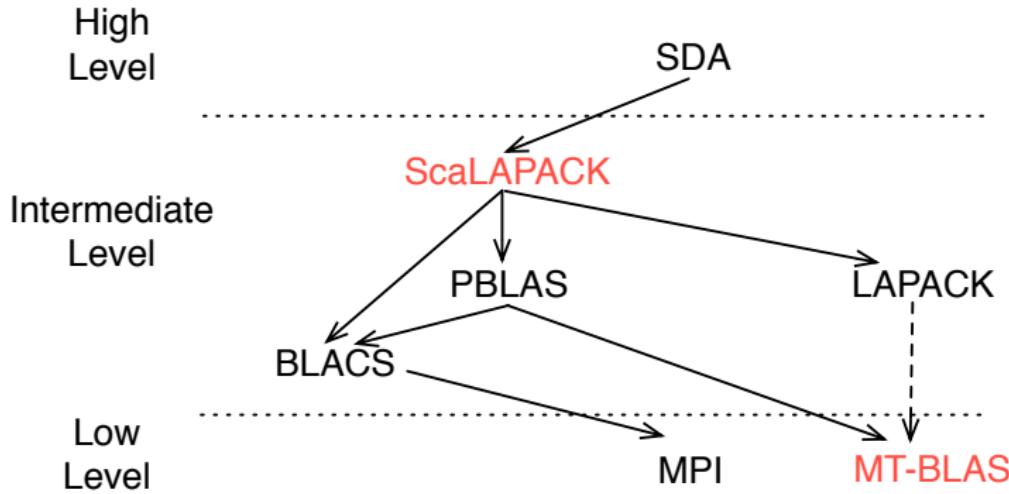
## Comments

- Can be used on shared and distributed memory multiprocessors.
- Similar performance than MT algorithm with few processors.  
Better with small problems.
- Improves MP and MT algorithms with 16 processors. Possible causes.
  - Exploits two different levels of parallelism: processes and threads.
  - Increases the locality of MT algorithm  $\Rightarrow$  Decreases data transference cost.

Could we take combine two different levels of parallelism? Last try.

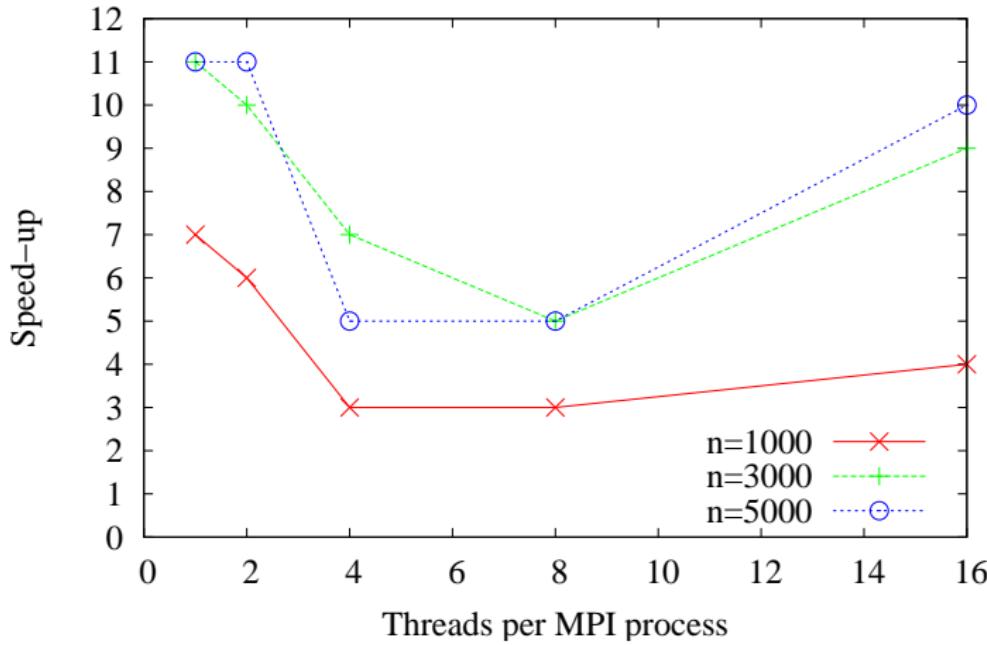


# Multiprocess hybrid algorithm (HBMP)



# Experimental results

Performance of the HBMP algorithm. 16 processors



## Comments

- The exact behaviour depends on the size of the problem.
- The hybrid algorithm (HBMP) is always improved by one of the "pure" algorithms: MT or MP algorithm.



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- We have applied five different parallel strategies to solve the Rational Matrix Equation.
- The algorithms exploit and combine different levels of parallelism and apply different programming models.
- Studying the dependencies of the algorithm and applying parallelism to the highest level pays off.
- The best results are obtained with a hybrid algorithm that combines processes with threads.
- The speed-up scales almost linearly with the number of processors with large problems.
- The algorithms use standard libraries ensuring the portability of the codes.