

Adaptive Precision Solvers for Sparse Linear Systems



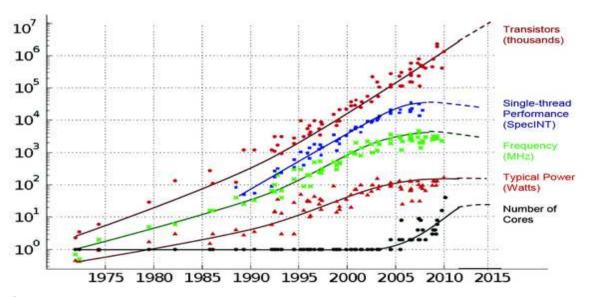
HARTWIG ANZT JACK DONGARRA



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Dennard's scaling vs Moore's Law

35 YEARS OF MICROPROCESSOR TREND DATA



Original data collected and plotted by M. Horowitz, F. Labonte, O. Shacham, K. Olukotun, L. Hammond and C. Batten Dotted line extrapolations by C. Moore



IBM Power8 (Q3'15) 22 nm 3.12 GHz

TDP 190-200 W

12 cores/96 threads



Intel Xeon E5-4669 v3 (Q2'15)

22 nm

2.1 GHz

TDP 135 W

18 cores/36 threads



- 5 nm in about 7-9 years:
 - 2.4*x* faster
 - 10x more transistors
 - ... but only 10% simultaneously active



Dark silicon (utilization wall) and specialization!





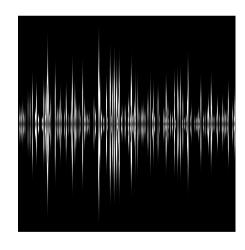








Approximate Computing: energy vs accuracy (or reliability)

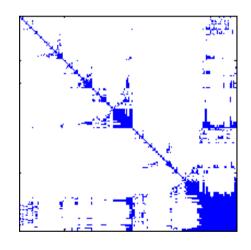


Signal & video processing

Probabilistic inference Service profiling Monte Carlo simulation Machine learning



Approximate Computing: energy vs accuracy (or reliability)



Numerical Linear Algebra?

- Tiny errors can rapidly aggregate
- Double precision is the standard



OUTLINE

- Jacobi solver for sparse linear systems
- Mantissa-adaptive Jacobi
- Experimental evaluation
- Concluding remarks



Stationary Methods: Jacobi

• Given Ax = b

$$x^{\{k\}} := D^{-1} \left(b - (A - D) x^{\{k-1\}} \right)$$

= $D^{-1} b + M x^{\{k-1\}}, \quad k = 1, 2, ...,$

with D=diag(A), and $x^{\{0\}}$ a starting solution guess

- Linear convergence provided spectral radius of M < 1
- Components of $x^{\{k\}}$ can be computed in parallel
- Alternative to exact triangular solves in approximate ILU preconditioning



- Basic idea:
 - Operate in low ("cheap") precision and gradually increase as needed
 - Currently {32,64}-bit precision
 - NVIDIA "Pascal" GPUs: {16,32,64}-bit precision?
 - FPGAs: custom







- Basic idea:
 - Operate in low ("cheap") precision and gradually increase as needed
 - Already explored for Jacobi

"On the potential of significance-driven execution for energy-aware HPC"
P. Gschwandtner, C. Chalios, D. S. Nikolopoulos, H. Vandierendonck, T. Fahringer
Computer Science – Research and Development, 2015

- ...but do it with a fine granularity (component-wise)
- ...and apply a cheap criterion to detect when to increase precision



- How?
 - Component-wise contraction property:

$$\left| x_i^{\{k\}} - x_i^{\{k-1\}} \right| \le \theta_i \left| x_i^{\{k-1\}} - x_i^{\{k-2\}} \right| \le \theta_i^2 \left| x_i^{\{k-2\}} - x_i^{\{k-3\}} \right| \dots$$



- How?
 - Component-wise contraction rate is constant:

$$c_i^{\{k\}} := \frac{z_i^{\{k-1\}}}{z_i^{\{k\}}} = \frac{\left|x_i^{\{k-1\}} - x_i^{\{k-2\}}\right|}{\left|x_i^{\{k\}} - x_i^{\{k-1\}}\right|}, \quad k \ge 2, \qquad c_i^{\{2\}} = c_i^{\{3\}} = c_4^{\{i\}} = \ldots = c_i$$

Exploding ratio $z^{\{k-1\}}/z^{\{k\}}$ due to small $z^{\{k\}}$ indicates convergence of the component in the current precision



- Practicalities:
 - Reduce the cost/periodicity of the test: $z^{\{k-\Phi\}}/z^{\{k\}}$
 - Take into account rounding errors:

$$\left|\frac{z_i^{\{k-\phi\}}}{z_i^{\{k\}}} - c_i^{\phi}\right| > \tilde{\delta}$$

determines that an extension is necessary

Avoid stagnation by setting

$$\tilde{\delta} := \delta \cdot (c_i^{\phi} - 1)$$

for some user-defined $0 < \delta < 1$



- Practicalities:
 - ullet δ governs how quickly the mantissa is extended:
 - \circ Faster/slower as $\delta \to 0/1$
 - Control the magnitude (gradient) of the increase: γ bits
 - First three iterations in full precision, to estimate component-wise contraction rate



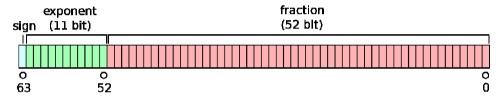
- Connection between fault tolerance and AC
 - Fault tolerance: obtain "exact" solution in presence of errors
 - AC: operating with low precision can be viewed as errors in the part of the mantissa that is chopped
 - Deviation from the contraction property: errors (fault tolerance) or convergence in current precisión (AC)

"Tuning iterative solvers for fault resilience" H. Anzt, J. Dongarra, E. S. Quintana-Ortí ScalA'2015 (Tomorrow!)



Experiments

Matlab (R2014a) and IEEE 754 double precision (64 bits):



- Solvers:
 - Full precision (52-bit mantissa)
 - 8-bit precision (8-bit mantissa)
 - Adaptive precision, starting with 8 bits
- Tune δ (tolerance threshold), γ (bit extension gradient), Φ (periodicity)



Experiments

- 27-pt stencil discretization of 3D Laplace using *16x16x16* mesh:
 - Symmetric matrix of order 4,096 with 97K nonzeros, well conditioned
 - Starting guess $x^{\{0\}} = 0, b = [1, 1, ..., 1]$
 - Convergence stopping criterion

$$\mathcal{R}\left(x^{\{k\}}\right) = \frac{\|b - Ax^{\{k\}}\|_2}{\|b - Ax^{\{0\}}\|_2}$$



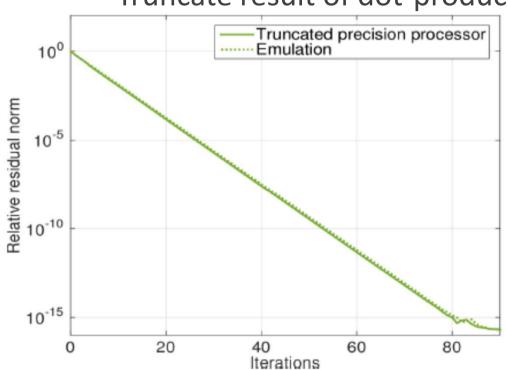
Experiments

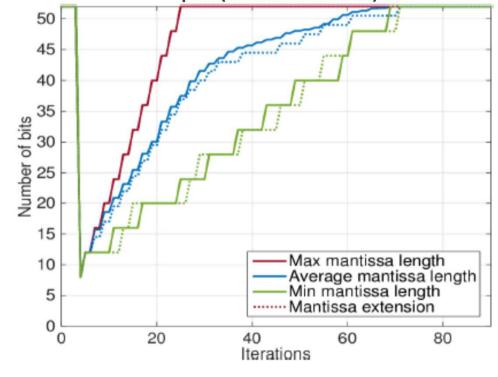
- Intel Xeon E5-2670 CPU. Emulate adaptive precision in hardware. At each iteration:
 - Truncate input data to desired precision
 - Compute (in 64-bit arithmetic) $y_i := A(i,:) \cdot x^{\{k-1\}}$ and truncate (to desired precision)
 - Compute $p_i := (b_i y_i) / A(i,i)$ and truncate
 - Compute next iterate $x^{\{k\}} := x^{\{k-1\}} + p_i$ and truncate



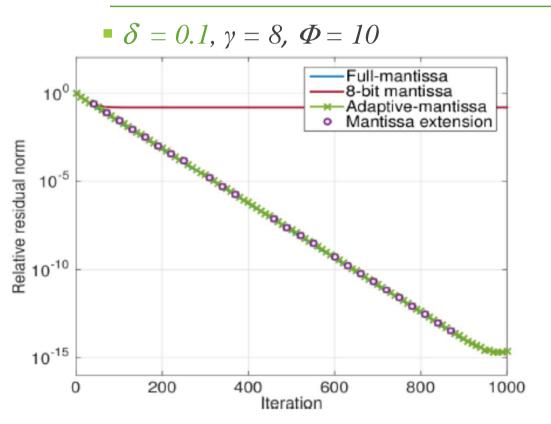
Experiments: validate emulation mode

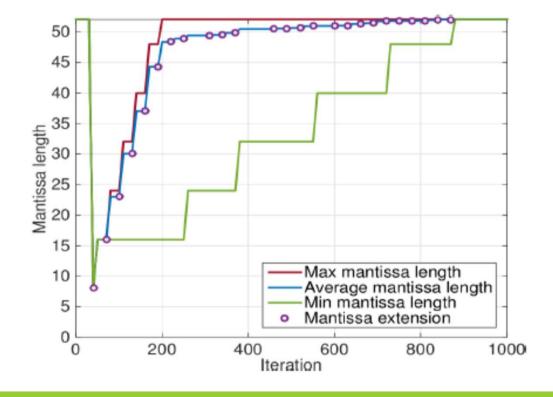
• Truncate result of dot-product vs truncate all flops (4x4x4 case)



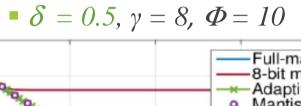


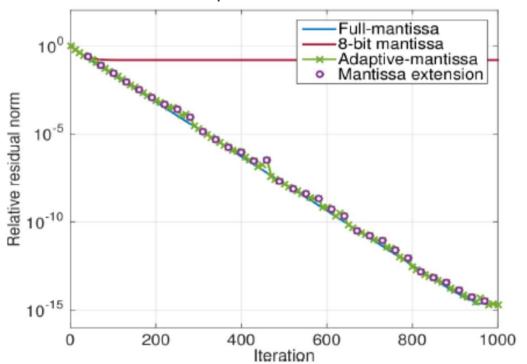


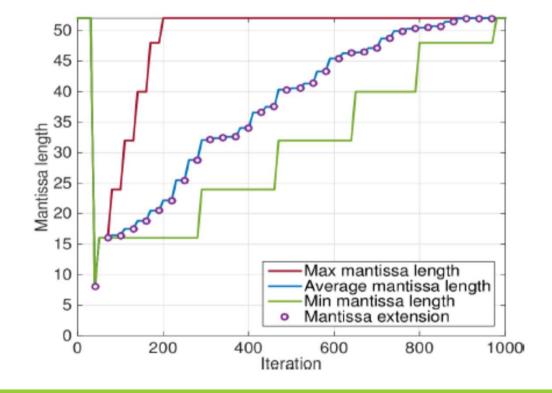




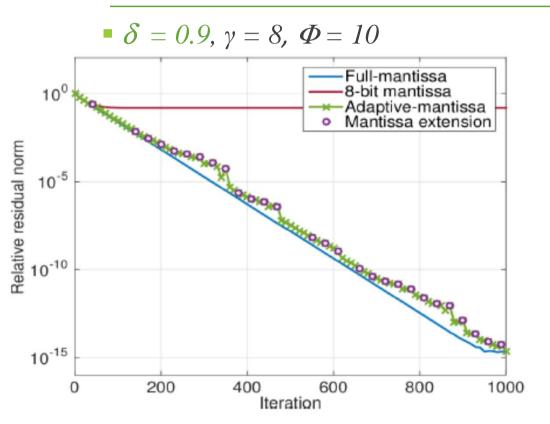


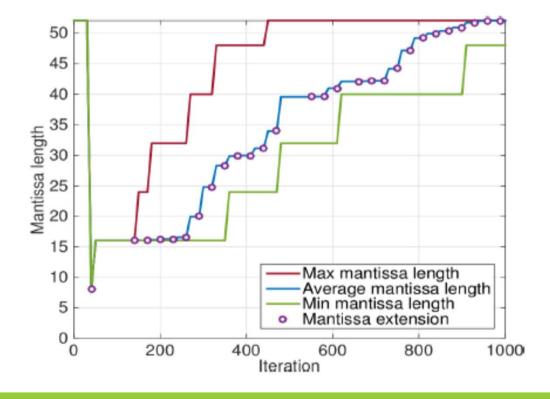












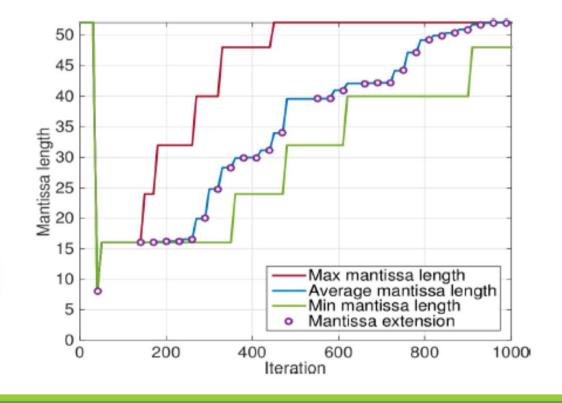


- Quantify computational cost/savings
 - For iteration k

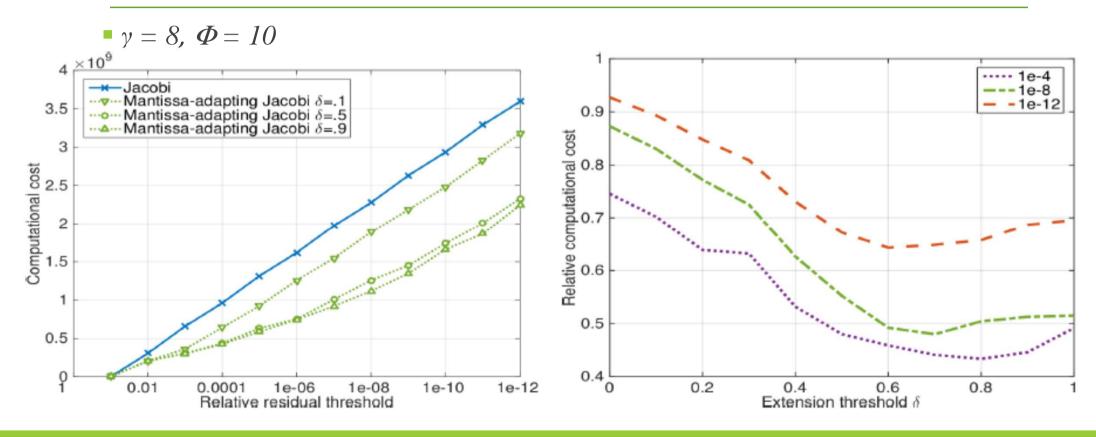
$$\mathscr{C}^{\{k\}} := \sum_{i=1}^n nnz_i(A) \cdot \mathscr{L}_i^{\{k\}}$$

For the full solve

$$\mathscr{C}^{[1,\tilde{k}]} := \sum_{k=1}^{\tilde{k}} \mathscr{C}^{\{k\}} = \sum_{k=1}^{\tilde{k}} \left(\sum_{i=1}^{n} nnz_i \cdot \mathscr{L}_i^{\{k\}} \right)$$



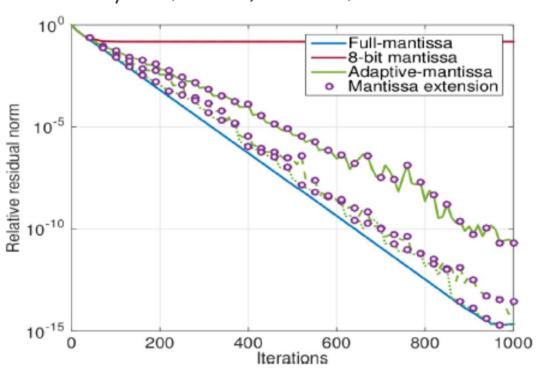


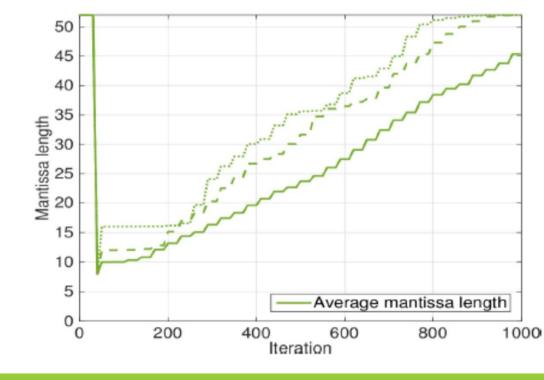




Experiments: select γ

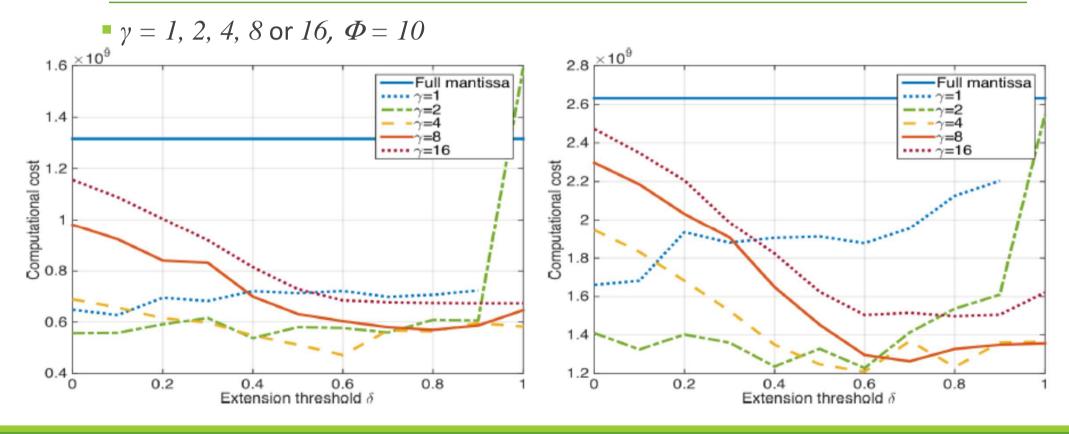
•
$$\gamma = 2$$
, 4 or 8, $\Phi = 10$, $\delta = 0.8$





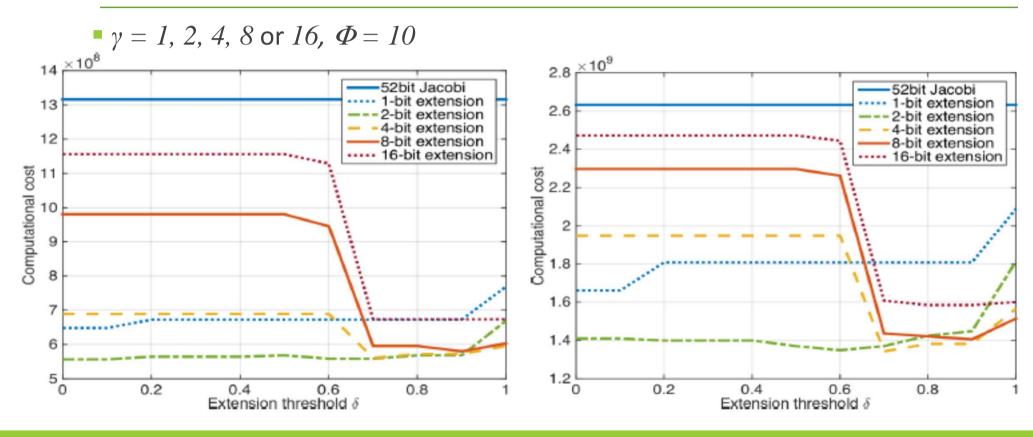


Experiments: select γ





Experiments: vector-wise adaptive





Experiments: Approx. triangular solves in ILU preconditioner

		Optimal			Number of iterations					Computational cost of adaptive-mantissa				
		configuration			Adaptive-mantissa Jacobi/full-mantissa Jacobi					w.r.t. full-mantissa Jacobi				
		φ	γ	δ	10^{-2}	10^{-4}	10^{-6}	10^{-8}	10^{-10}	10^{-2}	10^{-4}	10^{-6}	10^{-8}	10^{-10}
СНР	L	1	16	0.5	13/11	23/20	31/29	40/38	48/46	0.69	0.82	0.82	0.85	0.88
	\boldsymbol{U}	1	16	0.4	12/11	22/20	29/28	37/37	45/46	0.60	0.76	0.79	0.81	0.82
DC	L	1	32	0.8	12/8	13/10	15/11	17/13	19/13	1.31	1.13	1.21	1.17	1.34
	\boldsymbol{U}	1	32	0.8	14/8	15/10	16/11	18/12	19/13	1.59	1.35	1.31	1.37	1.34
LAP	\boldsymbol{L}	1	32	0.5	11/11	21/20	30/29	39/38	48/48	0.59	0.74	0.77	0.79	0.79
	\boldsymbol{U}	1	8	0.3	11/11	20/20	28/29	37/38	47/48	0.43	0.40	0.43	0.51	0.59
STO	L	1	16	0.7	10/9	21/20	31/28	38/34	46/44	0.74	0.87	0.98	1.01	0.96
	\boldsymbol{U}	1	8	0.1	8/9	20/20	28/28	35/36	45/45	0.51	0.65	0.75	0.78	0.85
VEN	L	1	16	0.9	25/24	43/42	60/59	74/73	86/85	0.82	0.89	0.92	0.94	0.95
	\boldsymbol{U}	1	8	0.6	17/17	36/35	53/52	68/68	82/82	0.54	0.77	0.85	0.87	0.89



Conclusions

- Careful exploitation of the component-wise contraction property of Jacobi iteration:
 - Monitor deviations from the expected convergence rate
 - Account for rounding error, but avoid stagnation
 - Cheap and periodic test for extension
- Potential savings of up to 60% for 3D Laplace benchmark
- Link with fault tolerance

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